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3.	ADDITIONAL FEES							
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121	30	221	23	or cover sheet				
139	130	139	130	Non-English specification				
147	2,520	147	2,520	For filing a request for reexamination				
112	2,320 920*	112	2,520 920*	Requesting publication of SIR prior to				
112	320	112	320	Examiner action				
113	1,840*	113	1,840*	Requesting publication of SIR after				
, 13	1,040	113	1,040	Examiner action				
115	110	215	55	Extension for response within first month				
116	380	216	190	Extension for response within second month				
117	870	217	435	Extension for response within second month				
118	1,360	217	680	Extension for response within fourth month				
128	1,850	228	925	Extension for response within fifth month				
119	300	219	150	Notice of Appeal				
120	300	220	150	Filing a brief in support of an appeal				
121	260	221	130	Request for oral hearing				
138	1,510	138	1,510	Petition to institute a public use proceeding				
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140	110	240	ออ	application				
141	1,210	241	605	Petition to revive unintentionally				
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142	1,210	242	605	Utility issue fee (or reissue)				
143	430	243	215	Design issue fee				
144	580	244	290	Plant issue fee				
122	130	122	130	Petitions to the Commissioner				
123	50	123	50	Petitions related to provisional applications				
126	240	126	240	Submission of Information Disclosure Stmt				
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### APPLICATION FOR UNITED STATES LETTERS OF PATENTS

#### **FOR**

# SYSTEM AND METHOD FOR SAFEGUARDING DATA BETWEEN A DEVICE DRIVER AND A DEVICE

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## SYSTEM AND METHOD FOR SAFEGUARDING DATA BETWEEN A DEVICE DRIVER AND A DEVICE

#### BACKGROUND OF THE INVENTION

#### 5 Field of the Invention

The present invention relates to data encryption. More specifically, the present invention relates to safeguarding the transfer of data within a device.

## Background Information

With the proliferation of computers and networks, the amount and availability of digitized data available for viewing and listening has grown. However, with this growth in the amount and availability of information, content providers have desired greater protection of the data from unauthorized use.

In order to protect data from unauthorized use, conventional data protection techniques, such as, for example, data encryption, have been used to protect data as it is being transferred over a network or between devices. Content providers use a number of well known encryption techniques to encrypt sensitive data before transmission from one device, such as, for example, a satellite receiving dish, to a second device, such as, for example, a computer or set-top box.

Different conventional types of encryption techniques are used depending upon the source device of the data and the type of data bus being used for the transmission from one device to another. For example, data transmitted from a Digital Video Disk (DVD) player to a computer uses Content Scrambling System (CSS) encryption, and data transmitted over an IEEE 1394 bus use Digital Transmission Content Protection (DTCP). Data transmitted over other bus systems

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use a number of other encryption techniques. In order to decrypt the data as it is received, devices need to be able to decrypt data using the variety of techniques that are used to encrypt the data. Thus, a device that receives both CSS and DTCP encrypted data needs to know the techniques for decrypting both types of encrypted data.

The various encryption techniques employed only protect the data during transmission. Once the data is received, it must be decrypted in order for the receiving device to be able to process the data. Once the data is decrypted within the receiving device, the data is susceptible to unauthorized access and manipulation.

Moreover, these conventional systems do not protect the data inside an open architecture device, such as a personal computer. Conventional systems do not control what applications access the incoming data-stream, nor allow those applications to access the incoming data stream without being aware of the data originator outside the device.

#### SUMMARY OF THE INVENTION

According to one aspect of the invention, a machine readable medium provides instructions which when executed by at least one processor, cause the processor to perform operations. The operations include encrypting a payload of a data-stream data block with at least one key before transmitting the data-stream from a first system to a second system, replacing a portion of the payload with a tag that identifies at least one decrypting key to the first system before transmitting the

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data-stream from the first system to the second system, and setting a flag in a header of the data block that indicates that the payload has the tag before transmitting the data-stream from the first system to the second system.

## 5 BRIEF DESCRIPTION OF THE DRAWINGS

The present invention will be described by way of exemplary embodiments, but not limitations, illustrated in the accompanying drawings. Identical numerals indicate the same elements throughout the figures.

Figure 1 is one embodiment for a data safeguarding system block diagram;

Figure 2 is one embodiment for an architecture of a data safeguarding system block diagram;

Figure 3 is another embodiment for an architecture of a data safeguarding system block diagram;

Figure 4 illustrates an exemplary architecture of a data safeguarding system, such as that shown in Figure 2;

Figure 5 is one embodiment for a protected content exchange (PCX) module of Figure 2 block diagram;

Figure 6a is one embodiment for an encrypted data stream block diagram;

Figure 6b is one embodiment for a PCX replacement block diagram;

Figure 7 is one embodiment for a shared buffer block diagram;

Figure 8 is one embodiment for a PCX resync block block diagram;

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Figure **9** is a flow diagram of one embodiment for safeguarding protocol specific data within a device;

Figure **10** is a flow diagram of one embodiment for decrypting PCX encrypted data by a decoding device;

Figure **11** is a flow diagram of one embodiment for creating a PCX resync block;

Figure **12** is a flow diagram of one embodiment for decrypting a PCX resync block;

Figure **13** is one embodiment for an information synchronizing system block diagram.

Figure **14** is one embodiment of a system block diagram showing the functional connection between a PCX module and an application decoder for transferring a data-stream to a decoder application when they are separate physical devices.

Figure **15** is an exemplary computer system that is related to the use of the present invention, according to an embodiment.

Figure **16** is one embodiment of a system block diagram showing the functional connection between a PCX module and an application decoder when they access a shared memory device.

Figure **17** is one embodiment of a system block diagram of a shared memory device safeguarding system.

Figure **18** is a flow diagram of one embodiment for transferring a single datastream and decryption keys to an application decoder.

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### **DETAILED DESCRIPTION**

In the following description, various aspects and details of the present invention will be described. However, it will be apparent to those skilled in the art that the present invention may be practiced with only some or all aspects of the present invention. For purposes of explanation, specific numbers, materials and configurations are set forth in order to provide a thorough understanding of the present invention. However, it will also be apparent to one skilled in the art that the present invention may be practiced without the specific aspects and details. In other instances, well known features are omitted or simplified in order not to obscure the present invention.

Some portions of the descriptions that follow are presented in terms of algorithms and symbolic representations of operations on data bits within a computer memory. These algorithmic descriptions and representations are the means used by those skilled in the data processing arts to most effectively convey the substance of their work to others skilled in the art. An algorithm is here, and generally, conceived to be a self-consistent sequence of steps leading to a desired result. The steps are those requiring physical manipulations of physical quantities. Usually, though not necessarily, these quantities take the form of electrical or magnetic signals capable of being stored, transferred, combined, compared, and otherwise manipulated. It has proven convenient at times, principally for reasons of common usage, to refer to these signals as bits, values, elements, symbols, characters, terms, numbers, or the like.

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It should be borne in mind, however, that all of these and similar terms are to be associated with the appropriate physical quantities and are merely convenient labels applied to these quantities. Unless specifically stated otherwise as apparent from the following discussion, it is appreciated that throughout the description, discussions utilizing terms such as "processing" or "computing" or "calculating" or "determining" or "displaying" or the like, refer to the action and processes of a circuit that can include a programmed computer system, or similar electronic computing device. A computer system manipulates and transforms data represented as physical (electronic) quantities within the computer system's registers and memories into other data similarly represented as physical quantities within the computer system memories or registers or other such information storage, transmission or display devices.

The present invention also relates to apparatus including circuits for performing the operations herein. This apparatus may be specially constructed for the required purposes, or it may include a general purpose computer selectively activated or reconfigured by a computer program stored in the computer. Such a computer program may be stored in a computer readable storage medium. A machine readable medium includes any mechanism that provides (i.e. stores and/or transmits) information in a form readable by a machine such as a computer. For example, a machine readable medium includes, and is not limited to, read only memory (ROM); random access memory (RAM); magnetic disk storage media; optical storage media; flash memory devices; electrical, optical, acoustical or other

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form of propogated signals (such as carrier waves, infrared signals, digital signals, and so forth)., or any type of media suitable for storing electronic instructions.

The algorithms and displays presented herein are not inherently related to any particular computer or other apparatus. Various general purpose systems may be used with programs in accordance with the teachings herein, or it may prove convenient to construct more specialized apparatus to perform the required method steps. The required structure for a variety of these systems will appear from the description below. In addition, the present invention is not described with reference to any particular programming language. It will be appreciated that a variety of programming languages may be used to implement the teachings of the invention as described herein.

Various operations will be described as multiple discrete steps performed in turn in a manner that is most helpful in understanding the present invention, however, the order of description should not be construed as to imply that these operations are necessarily order dependent, in particular, the order the steps are presented. Furthermore, the phrase "in one embodiment" will be used repeatedly, however the phrase does not necessarily refer to the same embodiment, although it may.

Figure 1 is a block diagram of one embodiment for a data safeguarding system 100. Data safeguarding system 100 includes data safeguarding device 104, protocol specific input devices 110 and protocol specific buses 120. Data safeguarding device 104 includes decoding devices 102, and a protected content exchange (PCX) module whose preferred embodiment includes a memory 108, and

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a CPU **115** that executes programmed instructions stored in a memory **108**. PCX module **106** includes a number of protocol specific exchange modules **130**.

Protocol specific encrypted data is received over protocol specific bus 120 from protocol specific input devices 110. In the Figure 1 example, encrypted data may be received over a 1394 DTCP bus from a number of input devices 110 such as a satellite dish or video recorder (VCR). Any of a number of protocol specific buses 120 may be connected to data safeguarding device 104 including, for example, a USB bus, a PCI bus, and a DVD bus. Once the encrypted data is received by data safeguarding device 104, CPU 115 directs the input to PCX module 106. Within PCX module 106, the appropriate protocol specific exchange module 130 is used to decrypt the encrypted input data stream. For example, if IEEE 1394 DTCP bus encrypted data is received, a DTCP exchange module 130 would be used to decrypt the input data. Input data is received and is decrypted on a block-by-block basis.

Initially, PCX module **106** negotiates a content channel encryption key with protocol specific input device **110**. PCX module **106** then negotiates a PCX session key with the client decoding device **102**. Decoding device **102** is the client that, in one embodiment, originally requested the data from device **110**. Once the PCX session key is negotiated, PCX module **106** re-encrypts the payload of the protocol specific data using a randomly generated PCX content key and transfers the re-encrypted data (including header and payload) to the appropriate decoding device **102**. Once decoding device **102** receives the re-encrypted data, decoding device

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102 negotiates with the PCX module 106 to retrieve the PCX content key encrypted by the PCX session key. Once the appropriate PCX content is retrieved, decoding device 102 decrypts the payload data. Decoding device 102 then manipulates the unencrypted data. In one embodiment, decoding device 102 decodes the unencrypted data. For example, if MPEG data is requested by an MPEG decoder, the appropriate input device 110 sends the data over the bus 120 to data safeguarding device 104. CPU 115 executes the PCX module 106 which decrypts the MPEG input data stream using a content channel encryption key for the bus 120. The MPEG decoder and PCX module 106 negotiate a PCX session key. The payload MPEG data is re-encrypted with the randomly generated PCX content key and the re-encrypted data is sent to the MPEG decoder. PCX module 106 encrypts the PCX content key with the PCX session key. The MPEG decoder retrieves the encrypted PCX content key and decrypts the PCX content key with the PCX session key. In addition, the MPEG decoder uses the PCX content key to decrypt the payload data for playback. The MPEG decoder then retrieves the device key and decrypts the payload data for playback.

In one embodiment, data within system **100** is further protected from tampering or from unauthorized access by the use of a number of anti-tampering techniques such as, for example, self-modification of PCX module **106** code, the use of anti-debugging techniques, self-verification of PCX module **106** code, signature verification of PCX module **106** code, and other applicable anti-tampering techniques. The use of these anti-tampering techniques prevents unauthorized

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access or modification of PCX module **106** code which prevents the unauthorized access or modification of the data as it is being transferred through system **100**.

Figure 2 is a block diagram of one embodiment for an architecture of a data safeguarding system 100. Referring to Figure 2, encrypted protocol specific data is received over IEEE 1394 bus 220 and transferred to IEEE 1394 bus driver 210. Bus driver 210 then sends the protocol specific data to class driver 212. PCX module 106 intercepts the protocol specific data and decrypts the data with a content channel encryption key. The content channel encryption key has originally been negotiated between PCX module 106 and protocol specific input device 110 before transmission. Once the data is decrypted, PCX module 106 re-encrypts only the MPEG portion of the payload of the data with a randomly generated PCX content key and encrypts the PCX content key with the appropriate PCX session key. This is repeated for the AC3 portion of the payload with a different randomly generated key and a different PCX session key. PCX module 106 sends the re-encrypted data back to class driver 212. The re-encrypted data is transferred to a splitter 232 which splits the data between the various decoding devices. In the figure 2 example, the splitter 232 splits the IEEE 1394 re-encrypted data to AC3 device 216 and MPEG device 218. MPEG decoder 218 and AC3 decoder 216 receive the appropriate encrypted PCX content key. MPEG decoder 218 and AC3 decoder 216 decrypt their PCX content key with their PCX session key. MPEG device 218 and AC3 device 216 then decrypt the re-encrypted data for playback using the appropriate PCX content key.

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Thus, the data is protected from unwarranted hacking or copying within data safeguarding system 100. Within data safeguarding system 100, the transmission headers of the data are left decrypted while the payload of the data is re-encrypted by PCX module 106. Thus, the payload of the data is protected from unwarranted copying or hacking during transfer within system 100 while allowing untrusted components to access the portions of the data stream they need.

Figure 3 is a block diagram of another embodiment of an architecture of a data safeguarding system 100. Referring to Figure 3, protocol specific input device 110 initially negotiates a content channel encryption key with protocol specific registration engine 326. Protocol specific input device 110 transmits the encrypted protocol specific data via protocol specific bus 120 to bus driver 312. Bus driver 312 transfers the encrypted protocol specific data to device specific mini port driver 316 via protocol specific class driver 314. Protocol specific bus abstractor 320 abstracts the encrypted protocol specific data from device specific mini port driver 316. The extracted encrypted data is transferred to PCX module 106. Within PCX module 106, the encrypted protocol specific data is decrypted using protocol specific decryptor 322. Protocol specific decryptor 322 decrypts the protocol specific data one block at a time. Each block of data contains a transmission header portion and a payload. In one embodiment, both the transmission header and payload portions are encrypted during transmission from source device 110 to data safeguarding system **100**. In an alternate embodiment, only the payload may be encrypted.

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Depending on the specific data bus transmission protocol being used, protocol specific decryptor **322** decrypts either the entire data block or the payload only.

Each data bus transmission protocol requires a corresponding protocol specific decryptor 322. PCX negotiator 328 negotiates a PCX session key with the decoding device 102 that is the intended recipient of the protocol specific data.

Once a session key is negotiated, protected content exchange (PCX) encryptor 324 re-encrypts the payload portion of the data with a randomly generated PCX content key to produce re-encrypted data. PCX encryptor 324 transfers the re-encrypted data to protocol specific bus abstractor 320 which, in turn, transfers the re-encrypted data to device specific mini port driver 316. Device specific mini port driver 316 sends the PCX re-encrypted data to the upstream drivers and libraries 330 which in turn transfers the PCX re-encrypted data to splitter 232.

Splitter 232 reads the transmission header of each re-encrypted data block and transfers the data block to the decoding device 102 corresponding to the information contained within the transmission header. In addition, in one embodiment, splitter 232 removes the transmission headers from the data block. Within the data, data blocks are intermingled so that a variety of data blocks are received by splitter 232. Thus, a video block may be received, then an audio block, then another video block, and so forth. The splitter transfers the payload sections of the blocks to the corresponding decoding device as indicated by the transmission header. Once the re-encrypted payload data is received by a decoding device 102, decoding device 102 retrieves the encrypted PCX content key from PCX negotiator

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**328**. Decoding device **102** decrypts the content key using its PCX session key which was originally negotiated with PCX negotiator **328**. The unencrypted data is then consumed by decoding device **102**.

Figure 4 illustrates an exemplary architecture of safeguarding system 100. Referring to figure 4, protocol specific input device 110, such as a VCR, negotiates with a playback device such as MPEG decoder 435 to transmit a stream of encrypted data to MPEG decoder 435. Protocol specific input device 110 initiates the transmission of a stream of encrypted protocol specific data marked with the appropriate copy protection status (i.e., "copy-1-generation," "copy-never," or "nomore-copies"). The copy protection status is transmitted via the encryption mode indicator (EMI) bits within the transmission header of the data. If data requested by decoding device 102 (such as an MPEG decoder 435) is copy protected, protocol specific input device 110 may choose to transmit an empty data stream until at least one decoding device 102 has completed the appropriate authentication procedure required to access the content stream. Within data safeguarding system 100, protocol specific input device 110 negotiates authentication through PCX negotiator 328 and not directly with protocol specific input device 110. In the figure 4 example, VCR 110 negotiates authentication with DTCP registration engine 426. Once protocol specific input device (VCR) 110 and DTCP registration engine 426 have completed the required AKE procedure, a content channel encryption key may be exchanged between protocol specific input device 110 and DTCP registration engine 426. This content channel encryption key is used to encrypt the data by

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protocol specific input device **110** and decrypt the IEEE 1394 encrypted data by DTCP decryptor **422**.

Once the content channel encryption key is negotiated, IEEE 1394 encrypted data is transferred from protocol specific input device 110 via IEEE 1394 bus driver 210, to class driver 212 and eventually to device specific mini port driver 416. DTCP bus abstractor 420 abstracts the IEEE 1394 encrypted data from device specific mini port driver 416 and transfers the IEEE 1394 encrypted data to PCX module 106. The IEEE 1394 encrypted data is decrypted by DTCP decryptor 422 one block at a time using the content channel encryption key previously negotiated by DTCP registration engine 426. In the IEEE 1394 example, both the transmission headers and the payload are encrypted by protocol specific input device 110. Thus, DTCP decryptor 422 decrypts both the transmission header and payload portions of the IEEE 1394 encrypted data block.

If video decoder **438** has not previously registered with PCX module **106**, PCX negotiator **428** authenticates video decoder **438**. During authentication, video decoder **438** is registered with PCX negotiator **428** and video decoder **438** negotiates a key exchange with PCX negotiator **428**. The key exchange method between video decoder **438** and PCX negotiator **428** is similar to the key exchange method between decoding device **110** and DTCP registration engine **426** described above. Once a session key is negotiated between video decoder **438** and PCX negotiator **428**, PCX encryptor **424** encrypts the payload of the data blocks using a randomly generated PCX content key. The re-encrypted IEEE 1394 data blocks are

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transferred to DTCP bus abstracter 420 for transfer to device specific mini port driver 416. The re-encrypted IEEE 1394 data is transferred via WDM stream class driver 430 and WDM streaming library 432 to source filter 434. At source filter 434, re-encrypted IEEE 1394 data intended for MPEG decoder 435 is split off from the other IEEE 1394 data and transferred to MPEG decoder 435. The re-encrypted IEEE 1394 data is muxed as MPEG transport stream (TS) to MPEG TS splitter 436. MPEG TS splitter 436 splits the video and audio portions of the MPEG TS and removes the transmission headers. The video portion of the TS is transferred to video decoder 438. Video decoder 438 requests the PCX content key from PCX negotiator 428. PCX negotiator 428 encrypts the PCX content key with the appropriate PCX session key and transfers it to video decoder 438. Video decoder 438 decrypts the PCX content key using the previously negotiated PCX session key and used the content key to decrypt the video data. In addition, the video decoder 438 consumes the data. In a similar manner, audio decoder 440 receives the audio TS and decodes the audio TS with a device key retrieved from PCX negotiator 428.

In standard MPEG video, the audio and video blocks are interwoven together within the input data stream. In order to separate the data, the MPEG splitter **436** reads the transport stream headers. Within data safeguarding system **100**, MPEG decoder **435** only needs to use the PCX specific protocols in order to interact with PCX negotiator **428** and does not need to be able to use each individual data bus transmission protocol. PCX module **106** is able to translate the encrypted protocol specific data from any specific bus into PCX encrypted data that the MPEG decoder

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435 is able to understand and decode. Thus, the re-encryption of the protocol specific data by PCX module 106 is independent of any specific bus protocol used by system 100. Decoding devices 102 are independent of the command protocol of the specific bus. The bus abstractor 420 abstracts the DTCP status structure, encapsulates the status structure in the proper command protocol, and transmits the encapsulated protocols to the driver 416 and vice versa. In this manner, decoding devices 102 are capable of receiving encrypted data from any protocol specific bus 120 without negotiating the content channel encryption key with the input devices 110 or knowing the encryption protocol for the specific buses 120. As existing bus protocols change and new bus protocols are developed, PCX module 106 may be updated. However, decoding devices 102 only need to be able to talk with PCX module 106 and only need to be updated when the PCX module 106 negotiation protocols are updated.

PCX module **106** may be implemented in software or hardware. The PCX module **106** may be incorporated within RAM memory of a personal computer or may be contained within flash memory which is attached to a CPU or other data processing device. Thus, PCX module **106** is easily updated independent of decoding devices **102**.

Figure **5** is a block diagram of one embodiment for a protected content exchange (PCX module **106**). Referring to **Figure 5**, PCX module **106** contains protocol specific decryption modules **500**, PCX encryption modules **510**, protocol specific registration modules **520**, and PCX negotiation modules **530**. A protocol

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specific decryption module **500** may be maintained for each protocol specific bus connected to data safeguarding system **100**. Thus, PCX module **106** may contain decryption module **1** (**502**) through decryption module n (**504**). PCX module **106** may contain a number of PCX encryption modules **510**. Thus, PCX module **106** may contain PCX encryption module **1** (**512**) through PCX encryption module n (**514**) for the encryption of a number of devices. In an alternate embodiment, only one PCX encryption module **510** may be maintained.

PCX module **106** includes a number of registration modules **520** for the negotiation of content channel encryption keys with protocol specific input devices **110**. In one embodiment, PCX module **106** may contain registration module **1** (**522**) through registration module n (**524**) corresponding to each protocol specific bus connected to the system.

PCX module **106** contains PCX negotiation modules **530** which are utilized by data safeguarding system **100** to negotiate key exchanges with decoding devices **102**. In addition, the negotiation modules authenticate the decoding devices and maintain key synchronization between PCX module **106** and decoding devices **102**. In one embodiment, PCX module **106** includes from negotiation module **1** (**532**) through negotiation module n (**534**) corresponding to individual decoding device **102**.

Figure **6A** is a block diagram of one embodiment of an encrypted data stream **600**. Referring to Figure **6A**, encrypted data stream **600** contains a number of blocks of data, each block containing a transport header **602** and a payload **604**. In

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one embodiment, the payload **604** and the transport stream header **602** may be 188 bytes in length. Within the encrypted data stream **600**, each block of data may be for a different device **102**. For example, MPEG audio and video data may be interleaved within encrypted data streams **600**. In addition, MPEG audio and video data may be interleaved with AC3 and other data.

Referring now to figure 6B, in an embodiment of the present invention a PCX data block 606 sent from a PCX module 106 to an application decoder 102 includes both a header 608 portion and a payload 616 portion. The header 608 portion is generally conventional and includes conventional block characteristic information, and a flag 609 of the present invention that indicates whether the payload 616 of the block data contains a tag 610, or alternatively whether the payload contains a PCX encrypted data. In one embodiment, the header 608 is a packetized elementary stream (PES) header. The payload 616 portion of the present invention includes the tag 610 at a predetermined position that includes an identifier information that can be sent to the PCX module for accessing the decryption key(s) for the payload as well as preferably a portion of the payload replaced by the tag, disclosed presently. The tag preferably includes a stream identifier datum 612 for distinctly identifying the data stream, and a source datum 614 for distinctly identifying the stream source, enabling the application decoder 102 to transmit to a PCX module a message that requests the decryption keys and preferably the portion of the payload for the identified data-stream from a PCX module that can access the decryption keys and preferably portion of the payload. In a safeguarding system 104 in which a data-

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stream identifier unambiguously includes the data sufficient to access the decryption keys and preferably the portion of the payload, the tag should only include the data-stream identifier. In other systems, particularly those have a plural number of PCX modules, the tag should also include an additional datum such as the source datum 614. When the payload includes the tag 610, the encrypted data stream is modified to replace a portion of the payload that is the size of the tag, with the tag. Thus, the payload content data 616 of the present invention is an encrypted form of the conventional data block that has a smaller portion replaced by the tag 610. This shall be presented more in subsequent paragraphs with reference to figures 14 and 17.

Referring now to figure 14, the block diagram depicted includes the PCX module 106, and the decoders 102, that contain circuitry of the present invention. The preferred embodiment of the application decoder 102 and the PCX module 106 each include a processing unit that responds to program instructions of the present invention. Alternatively, as is well known to practitioners of the art, the circuitry does not require a processing unit and can be implemented as a fixed digital circuit without the configurable circuit advantages provided by a programmed processing unit.

The source device **110** transmits an exemplary two intertwined data-streams, a video data-stream and an audio data-stream, to a device specific driver stack **1410** of data safeguarding device **104** via a bus **1420a**. Each data-stream includes a sequence of data blocks, each data block having a conventional header and

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payload. The driver stack **1410** retransmits each data-stream to an appropriate PCX module **106**. The PCX module **106** includes at least one decryptor and protocol specific registration engine, and at least one PCX encryptor and PCX negotiator, described herein with reference to figures **3** and **4**. Each data-stream transmitted from the source device **110** is optionally encrypted. The data-stream payloads are each encrypted by a PCX module **106** before transmission to an application decoder **102**, or alternatively optionally encrypted by a PCX module **106** if an individual data stream was transmitted from a source device **110** encrypted, and subsequently decrypted, by the PCX module **106**, so as to distinctly encrypt the data within the data safeguarding device **104**.

The embodiment portrayed in figure 14 includes an application decoders

102a and 102b that are each a physically separate device from the PCX module

106. There are two separate data transmission channels connecting the PCX

module 106 to each physically separate application decoder 102a and 102b. One

of the separate data transmission channels transmits the data-stream from the PCX

module 106 to the application decoder. The other separate data transmission

channel transmits the non-data-stream data between the PCX module 106 and an

application decoder 102, so these transmissions do not impact other components

that access the data-stream transmission. In the embodiment portrayed in figure 14,
each channel is a separate physical transmission line.

The data-stream data transmission path includes the PCX module **106** that sends the exemplary intertwined data-stream to a driver stack **1410**. The driver

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described with reference to figure 17.

stack 1410 sends the data-stream to a splitter 1432, wherein each separate datastream is then separated and separately transmitted to an appropriate exemplary application decoder **102a** or **102b**. The video data-stream is routed to the exemplary video application decoder 102a, and the exemplary audio data-stream is routed to the exemplary audio application decoder 102b. The non-data-stream data transmission path between the PCX module 106 and the decoder 102a is exemplary bus 1460a, and between the PCX module 106 and the decoder 102b is exemplary bus 1460b, wherein buses 1460a and 1460b may be identical physical devices. The non-data-stream data includes the identifier necessary for the PCX module to access the data block decryptor keys and optional portion of the payload. The nondata-stream data preferably includes a data-stream identification datum and a source identification datum from the decoders 102a and 102b, and the encryption keys and the portion of a replaced payload from the PCX module 106. The preferred embodiment non-data-stream data additionally includes an authentication and key exchange (AKE) from the PCX module 106 to the exemplary application decoders 102a and 102b to enable a separately encrypted tag and the aforementioned encryption keys to be themselves encrypted, assuring the embodiment of an authorized and secure decoder(s) 102 in communication with the PCX 106 module and receiving the data-stream. The precise method of transmitting and receiving the data-streams, datum identifiers, and encryption keys, shall be

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Referring now to figure 15, a programmed processor embodiment of the PCX module 106 runs on a computer system that can include an exemplary unitary processor 1510 that processes data signals. The processor 1510 may be a complex instruction set computer (CISC) microprocessor, a reduced instruction set computing (RISC) microprocessor, a very long instruction word (VLIW) microprocessor, a processor implementing a combination of instruction sets, or other processor device. However, it is understood that the present invention may be implemented in a computer system having multiple processors. The processor 1510 is coupled to a CPU bus 1520, or other communication device for communicating information, that transmits data signals between processor 1510 and other components in the PCX module 106. The computer system includes a memory 1530, or other computer readable media that is commonly a random access memory (RAM) device or other dynamic storage device, that can be used to store temporary variables or other intermediate information during execution of instructions by processor 1510, and is coupled to the bus 1520. The PCX module 106 also includes a read only non-volatile memory such as a semiconductor Read Only Memory (ROM) device, and/or other static storage device **1540** coupled to bus 1520 for storing static information and instructions for processor 1510. Data storage device 1550 is another computer readable medium coupled to bus 1520 for storing information and instructions, and can be such exemplary computer readable media as magnetic disk, and/or an optical disk and corresponding drives. Display 1560 is coupled to bus 1520 for displaying data generated by the processor 1510, and

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mouse **1570**, or other exemplary selecting or pointing device, and keyboard **1580**, each couple to the bus **1520**.

Referring to figure 16, a PCX module 106 includes a programmed processing device 1605 that accesses a memory unit 1615 for transmission of the encrypted data stream to that memory unit 1615, and for transmission of the keys(s) and tag data. The system includes the exemplary application decoders 1610, embodied by an exemplary video data application decoder 1610a and an exemplary audio data application decoder 1610b. The application decoders 1610 each access the memory unit 1615 for the encrypted data stream. The tag data is read by the decoders 1610, and sent back to the memory unit 1615, for access by the PCX computing device 1605, and a placement of the relevant key(s) and portion of the payload into a memory location that a decoder 1610a or 1610b accesses for a read of the key(s) and the replaced portion of payload data. Alternatively, the PCX computing device 1605 can store the key(s) and payload portion in the memory unit for a direct read by an application decoder 1610 according to the content of the transmitted tag data. In another embodiment, as disclosed herein, the application decoder(s) 1610 and the PCX computing device can be embodied by a unitary computing device that executes both program instructions for the application decoder(s), and the PCX module.

Referring to figure 17, a preferred embodiment block diagram depicted includes the PCX module 1706, the decoders 1702, and the driver stack 1710 that contain circuitry of the present invention. As formerly described with reference to

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figure 14, the source device 110 transmits an exemplary two intertwined datastreams to a device specific driver stack 1710 of data safeguarding device 104 via a
bus 1420a. The data safeguarding device 104 includes a shared memory 1715.

The driver stack 1710 moves each block to memory 1715 where it is written into a
buffer 1715a of the memory 1715, and sends to the PCX module 1706 a pointer to
the buffer 1715a for each block. The PCX module 1706 accesses each block
according to its memory pointer and distinctly encrypts the data within the
safeguarding device 104 as described with reference to figure 14.

The PCX module 1706 additionally replaces a portion of the payload with the tag, and marks a flag, as described with reference to figures 6b and 14, and as will be described with reference to figure 18. The memory 1715 includes a second buffer 1715b that both the exemplary decoders 1702a and 1702b and the PCX module 1706 write to and read from for transmission between them of non-data stream data described with reference to figure 14, and figure 18. The PCX module may also include a splitter circuit that places a pointer in the buffer 1715b identifying to the application decoders 1702 the data-streams directed to each separate exemplary application 1702a and 1702b, ort alternatively transmit that data over a separate physical line directly to the application decoders 1702 in a configuration that includes a pre-existing physical bus as depicted with reference to figure 14. The splitter circuit may be physically separate form the PCX module 1706 including a separate processor that may receive pointers directly from the driver stack 1710, and may write into a separate buffer in the memory 1710. In the embodiment herein

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portrayed. The interface between a decoder **1702a** and **1702b** and the buffer **1715a** is a first channel, and the interface between a decoder **1702a** and **1702b** and the buffer **1715b** is a second channel.

Referring now to figure 18, the method and circuit herein described applies to a system of a decoding application 102, portrayed with reference to both figure 14, wherein an exemplary video decoder 102a and audio decoder 102b, and a physically separate PCX module 106, in which a data stream is sent to the PCX module from a source device 110; and analogously to figure 17 as an exemplary video decoder 1702a and audio decoder 17102b, and a physically separate PCX module 1706; as well as a system implemented by a processing device that is both a PCX module and an application decoder(s). As has been described with reference to figure 14, the preferred circuit includes a programmed processing device, but alternatively can be implemented by digital circuitry that does not include a programmed processing device, or can be implemented alternatively by a programmed processing device in at least one application decoder and/or the PCX module, or a processing device that is embodied partially, but not completely, by a programmed processing device.

The data stream transmitted to the safeguarding system is alternatively unencrypted, or encrypted and has been decrypted by the PCX module as described herein. At block **1805**, the PCX module not necessarily but preferably performs an AKE procedure with each decoder to create a shared session key with each decoder. This session key will be used to encrypt the decryption keys before

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they are sent back to the decoder. Additionally this AKE will assure that the applications are authorized to access the PCX module encryption system. At block 1810, the PCX module encrypts the data block payload. The payload is encrypted using at least one key. At block 1815, the PCX source module stores a tag-sized portion of the encrypted payload for subsequent transmission to an application decoder. In the preferred embodiment, the entire payload is encrypted using the key(s). In the present invention, the stored portion can alternatively be encrypted separately with the key(s), or can be optionally left unencrypted. The payload in a following block shall be decrypted in accordance with the encryption characteristic of the stored portion.

At block **1820**, a tag is inserted into the payload in the place of the saved payload portion. The tag includes in the preferred embodiment both an identification of the data stream **612** and an identification of the data stream source **614**, the source identified because a safeguarding system may include more than one source circuit. The encryption keys and the saved portion of the payload are each referenced to the data-stream identifier. At block **1825**, a flag in the header is marked to indicate that the block contains a payload tag. At block **1830**, the data block is sent to the appropriate decoder **102** along the data-stream transmission channel described with reference to figure **14**, or alternatively described with reference to figure **17**. At block **1835**, the appropriate application decoder has received the data block from the splitter **1432** with reference to figure **14**. At block **1840** the application decoder that has received the data block reads the header flag

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position and at block **1845** determines whether the header flag is marked. If the header flag indicates that the payload does not contain a marked flag, control passes out of this flow. If the header flag indicates that the payload does contain a tag, control passes to block **1850** where the data stream identifier datum and the source datum are read and an identifier of each is sent back to each PCX module or alternatively, only the data stream identifier is sent back to the source module circuit identified by the source datum. In the embodiment in which the application decoder module, and the PCX module are physically separate devices, the identifier(s) are sent back to the PCX module along the separate channel as herein described.

At block 1855 the appropriate PCX module reads the data stream identifier. The proper application keys and portion of the payload are determined by reference to the data stream identifier. The second set of encryption key(s) and the stored portion of the payload that was replaced by the tag are transmitted to the target application decoder in accordance with the data stream identifier. In the embodiment in which the application decoder module and the PCX module are physically separate devices, the identifiers are sent back to the PCX module along the separate channel as herein described. At block 1860, the appropriate application decoder receives the decryption keys key(s) and the payload portion transmitted from the PCX module at block 1855, and decrypts the key(s) with the session key, replaces the payload portion from the tag position, and then decrypts the payload using the decrypted key(s).

Figure 7 is a block diagram of one embodiment for a shared buffer 700.

Shared buffer 700 includes a device specific header 710 and PCX resync blocks

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720. Device specific header 710 includes a header data portion 712 and PCX content key 714. In one embodiment, PCX resync blocks 720 contain from PCX resync block 1 (722) through PCX resync block n (726). Header data 712 identifies the decoding device 102 corresponding to the shared buffer 700. In one embodiment, each decoding device 102 corresponds to a unique shared buffer 700. In an alternate embodiment, all decoding device 102 use a single, shared buffer 700. Shared buffer 700 may be any applicable data structure such as, for example, an array, linked list, or other applicable data structure. PCX content key 714 is encrypted with the previously negotiated PCX session key and is the key that will be used to decrypt the payload.

Figure 8 is a block diagram of one embodiment for PCX resync block 720.

Referring to figure 8, PCX resync block 720 includes key delta tag 810, random initialization vector 815, and portion of the encrypted payload data 820. PCX resync block 720 is utilized for key synchronization as described below.

Figure **9** is a flow diagram of one embodiment for safeguarding protocol specific data within a device. Initially at processing block **905**, data safeguarding system **100** receives encrypted protocol specific data. The encrypted protocol specific data may be encrypted for any of a variety of data bus security protocols such as, but not limited to Digital Transmission Content Protection (DTCP), Content Scramble Systems (CSS), and Content Protection for Recordable Media (CPRM). The protocol specific data is received in processing blocks one block at a time.

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At processing block 910, the encrypted protocol specific data is translated into protected content exchange (PCX) re-encrypted data. The translation of the data includes decrypting the encrypted protocol specific data using a content channel encryption key to produce decrypted data. Once the data is decrypted, the payload of the decrypted data is re-encrypted using a PCX content key to produce PCX re-encrypted data. The content channel encryption key is negotiated by a protocol specific registration engine 326 with protocol specific input device 110 upon initiation of the transfer of protocol specific data from the protocol specific input device 110 to decoding device 102. Once protocol specific input device 110 and protocol specific registration engine 326 have completed the required AKE procedure, a content channel encryption key may be exchanged between protocol specific input device 110 and protocol specific registration engine 326. This content channel encryption key is used to encrypt the data by protocol specific input device 110 and decrypt the encrypted protocol specific data by protocol specific decryptor 322. The session key is negotiated between PCX negotiator 328 and decoding device 102.

After the data is re-encrypted, the re-encrypted data and the PCX content key encrypted by the PCX session key are transferred to the decoding device **102** at processing block **915**. In one embodiment, the re-encrypted data is split into a number of data streams which are transferred to appropriate decoding devices **102**.At processing block **920**, decoding device **102** decrypts the PCX content key

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and uses it to decrypt the re-encrypted data. The unencrypted data is further decoded by decoding device **102**.

Figure 10 is a flow diagram of one embodiment for decrypting re-encrypted data by decoding device 102. Referring to Figure 10, decoding device 102 receives re-encrypted data at processing block 1005. At processing block 1010, decoding device 102 retrieves the encrypted PCX content key from PCX negotiator 328. If decoding device 102 is not registered, PCX negotiator 328 registers the protocol device 102 and negotiates the PCX session key for the protocol device 102. At processing block 1015, decoding device 102 decrypts the re-encrypted data using the PCX content key.

Figure 11 is a flow diagram of one embodiment for creating a PCX resync block 720. Initially at processing block 1105, PCX module 106 receives protocol specific encrypted data. Next, at processing block 1110, PCX module 106 determines if a new resync point has been reached. If a new resync point has not been reached, processing continues at processing block 1130. If a new resync block has been reached, processing continues at block 1111. At processing block 1111, PCX module 106 determines if PCX content key needs to be generated. If no new PCX content key needs to be generated, processing block 1115. However, if a new PCX content key needs to be generated, processing continues at processing block 1112.

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At processing block **1112**, the new PCX content key is generated. PCX module **106** uses the existence of natural synchronization points within the original data stream to determine when to create a new PCX content key.

At processing block 1115, PCX module 106 generates PCX tag 610 that is a unique identification for the PCX resync block 720. In one embodiment, PCX tag 610 may be an array index value. In alternate embodiments, PCX tag 610 may be any suitable index value to the PCX resync block 720. At processing block 1120, PCX module 106 copies PCX flag 609, PCX tag 610, TSID 612, and PID 614 into the payload portion of the data stream and saves the original portion in location 820 in the resync block 720.

At processing block 1125, PCX module 106 updates PCX resync data 720. If the PCX content key being used to encrypt the payload is different from the PCX content key used on the previous block for the same decoding device 102, key delta tag 810 is incremented. Otherwise, key delta tag 810 is unchanged. In this manner, PCX content keys may be changed periodically during re-encryption of the data. This increases the security of the data within system 100. In one embodiment, PCX content key is changed on a fixed time interval or after a fixed number of PES headers 608 have been processed.

In order to increase the security of system 100, the PCX content key is altered on each PES header 608 change by using a random initialization vector as a seed value to modify the key. This allows splitter 232 to drop a data block without losing the ability to decrypt the remaining data in the input stream. In one

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embodiment, key delta tag **810** and random initialization vector **815** are not encrypted. PCX content key **714** is encrypted with the previously negotiated PCX session key.

At processing block **1130**, PCX module **106** encrypts the payload containing the resync data using the PCX content key.

Figure 12 is a flow diagram of one embodiment for decrypting a PCX resync block 720. Initially at processing block 1205, decoding device 102 receives a block of PCX encrypted data. At processing block 1210, decoding device 102 decrypts the payload and determines if the block of data is a resync block. If not, processing continues at step 1219. If the block of data is a resync block, processing continues at block 1211.

At processing block 1211, decoder 102 checks if key delta tag 810 changed.

Delta tag 810 indicates if PCX content key has changed. If so, at processing block

1213, decoding device 102 retrieves PCX content key 714 from shared buffer 700.

At processing block 1215, decoding device 102 extracts PCX tag 610 and performs a look-up of the resync block 720 within shared buffer 700. Decoding device 102 restores the original payload.

Decoding device **102** then decrypts the PCX content key using the previously negotiated PCX session key. At processing block **1218**, decoder **102** reinitializes the decryption cipher using the PCX content key and the random initialization vector **815**.

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At processing block **1219**, decoder **102** decrypts the payload using the decryption cipher. At processing block **1220**, the decoding device **102** decodes the payload of the unencrypted data for further processing (for example, playback by MPEG decoder).

The protocol specific data may contain copy control information (CCI) which allows the content owners to assign varying levels of priority for what can and can't be done with the data. The data may be "copy free" which means there is no restriction to copying the data. The other end of the spectrum is "copy never" which means that as soon as the AKE is negotiated, a device must render the data immediately. In this scheme, a device can not make any copies, can not save the data for later use, or anything similar. Thus, when a device receives the data, it is sent to the consumer, and then the data gets thrown away.

The other two schemes are "copy once" and "copy no more." If a device receives data that is marked as "copy once," the device may make a single copy of the data if the user chooses to do so. This scheme allows recording for later viewing. When a device receives data that is marked "copy once," the device may save it, but then once it is saved, when it is retrieved after saving, the device must mark the data as "copy no more."

In one embodiment, during transfer of data within system **100**, if the data is unencrypted, the CCI information is susceptible to interception and unauthorized change. Thus, if the data is marked "copy never" and the information is hacked, the data may be pirated within system **100**. The CCI information is contained within

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transmission header **602**. The transmission header **602** is not encrypted during transfer though system **100** and is susceptible to change.

Within system 100, the CCI information is built into the PCX content key. The CCI information retrieved from the data stream in transmission header 602 is used as part of the seed to generate the key. Thus, by combining the PCX content key with the control information before re-encryption, system 100 guarantees that any modification of the CCI information in the transmission header 602 will result in incorrect decryption of the protected data. During decryption of the re-encrypted data by decoding device 102, the CCI information is extracted from the transmission header 602 and combined with the PCX content key to create the decryption key.

The above method may be used to protect any information embedded within the transmission header **602**. Thus, information such as, for example, copy quality which may indicate the quality of audio a user is allowed to copy, how many times a device is allowed to copy this content, and similar information may be protected from change while the data is transferred within system **100**.

Figure 13 is a block diagram of one embodiment for an information synchronizing system 1500. Content exchange device 1510 is configured to receive fixed-size data 1505. Content exchange device 1510 is further configured to save a portion of the original payload of the fixed-size data 1505 in shared memory buffer 1540 and configured to save synchronization information together with the original portion in shared memory buffer 1540. In one embodiment, decryptor 1525 is configured to decrypt fixed-length data 1505 as it is received by content exchange

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device **1510**. Negotiator **1515** is configured to embed a tag to the appropriate synch block in shared memory buffer **1540** within a payload area of the fixed-size data **1505** to produce replacement data **1530**. In one embodiment, encryptor **1520** is configured to encrypt the payload of replacement data **1530** and configured to encrypt the original payload saved in shared memory buffer **1540**.

Decoding device **1535** is configured to extract the embedded tag from replacement data **1530** and to retrieve the original payload and synchronization information from shared memory buffer **1540** corresponding to replacement data **1530**.

In one embodiment, decoding device **1535** is contained within the same device as shared memory buffer **1540**. In an alternate embodiment, decoding device **1535** is a separate device from the device containing shared memory buffer **1540**.

While certain exemplary embodiments have been described and shown in the accompanying drawings, it is to be understood that these embodiments are merely illustrative of and not restrictive of the broad invention. The present invention is not limited top the specific constructions and arrangements shown and described, and alternative embodiments will become apparent to those skilled in the art to which the present invention pertains without departing from the scope of the present invention. The scope of the present invention is defined by the appended claims rather than the forgoing description. In the appended claims, a physical embodiment of each recited circuit limitation does not necessarily include completely separate

physical devices from another recited circuit limitation. An embodiment of each circuit may share at least one element with another circuit.

# **CLAIMS**

What is claimed is:

1	1. A machine readable medium that provides instructions, which when
2	executed by at least one processor, cause said processor to perform operations
3	comprising:
4	encrypting a payload of a data block of a data-stream with at least one key,
5	before transmitting the data-stream from a first system to a second system;
6	replacing a portion of said payload with a tag that identifies an at least one
7	decrypting key to said first system, before said transmitting; and
8	setting a flag in a header of the data block that indicates that said payload
9	has said tag, before said transmitting.
1	2. The medium defined in claim 1 wherein said encrypting includes encrypting
2	said portion of said payload.

- The medium defined in claim 1 wherein said tag includes one of: 3. 1
- a data-stream identifier, and 2
- a data-stream identifier and a source, said source characterized by at 3 least one of a source of said keys, and a source of said keys and a source of 4 said portion of said payload. 5
- The medium defined in claim 1 wherein said operations further include 1 4.

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- receiving a transmission from said second system that includes data 2 3 indicating said tag; and
- sending one of said keys, and said keys and said portion of said payload, to 4 5 said second system based on said transmission.
- The medium defined in claim 1 wherein said operations further include before 5. 1 2 setting said flag and encrypting said payload; said first system
- 3 setting said flag in said header,
- 4 encrypting said payload, and
- receiving a stream of data from a third system wherein said data-stream is 5 6 based on said stream of data.
- 6. A machine readable medium that provides instructions, which when executed 2 by at least one processor, cause said processor to perform operations comprising:
  - after a fixed-length data block of a data-stream, the data block having both a payload including an encrypted data portion and at least one tag bits, and a header, is received by a second system, reading a flag in the header indicating that the data block has the tag bits;
  - if the flag indicates that the data block has the tag bits, reading at least one bit identifying the data-stream in the tag bits;
- 9 sending a datum from the second system to a transmitting first system 10 indicating an identification of the read data-stream based on the at least one bit;

the second system receiving from the first system a definition of a decrypting keys for the data-stream based on the datum sent from the second system to the first system; and

decrypting the data block in the second system based on the decrypting keys received by the second system.

- 1 7. The medium defined in claim 6 further including the second system receiving
- 2 from the first system the portion of the payload based on the datum sent from the
- 3 second system to the first system.
- 1 8. The medium defined in claim 6 further including the second system replacing
- 2 the at least one tag bits in the payload with the portion of the payload, and if the
- 3 portion of the payload is encrypted the decrypting includes decrypting the portion of
- 4 the payload.
- 1 9. A method comprising:
- a sending system replacing a portion of a data block payload with at least one
- 3 tag bits that identify an at least one decrypting key;
- 4 said sending system setting a flag in a header of said data block that
- 5 indicates at least one of said payload is encrypted and said payload includes said
- 6 tag;
- 7 said sending system encrypting said payload with at least one key; and
- 8 said sending system transmitting said data block to a receiving system after
- 9 said setting a flag, said encrypting, and said replacing.

- 1 10. The method defined in claim 9 wherein said encrypting includes encrypting
- 2 said payload portion.
- 1 11. The method defined in claim 9 further including said sending system
- 2 transferring a first data characterized by at least one of:
- 3 said at least one key to said receiving system; and
- said at least one key and said payload portion to said receiving system.
- 1 12. The method defined in claim 11 wherein said sending system transmitting
- 2 said first data is based upon said receiving system transmitting to said sending
- 3 system said tag bits.
- 1 13. The method defined in claim 12 further including one of:
- 2 said sending system transmitting said payload portion to said receiving
- 3 system based upon said receiving system transmitting to said sending system said
- 4 tag bits; and said receiving system replacing said tag bits with said payload portion
- 5 in response to receiving said payload portion from said sending system, and
- 6 wherein said encrypting includes encrypting said payload portion, and said
- 7 decrypting includes decrypting said payload portion; and
- 8 said sending system transmitting said payload portion to said receiving
- 9 system based upon said receiving system transmitting to said sending system a first
- datum that identifies a data-stream that includes said data block, and said receiving

- 11 system replacing said payload portion in response to receiving said payload portion
- 12 from said sending system.
- 1 14. The method defined in claim 9 wherein said transmitting occurs via a shared
- 2 memory unit.
- 1 15. The method defined in claim 9 wherein
- 2 said sending system and said receiving system are separate physical
- 3 devices;
- 4 said transmitting of said data block occurs on a first channel; and
- 5 transmitting of non-data block data including at least one of said key from
- 6 said sending system to said receiving system, said payload portion from said
- 7 sending system to said receiving system, and a datum that identifies a data-stream
- 8 that includes said data block, occurs on at least one separate second channel.
- 1 16. The method defined in claim 9 wherein said tag bits further identify a source
- 2 of said keys in said sending system.
- 1 17. A method comprising:
- 2 a receiving system of an encrpted data block that has a payload and a
- 3 header reading a set flag in a header of said data block;
- 4 said receiving system reading at least one tag bit in a payload portion of said
- 5 data block in response to said reading said set flag;

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6	said receiving system sending a first datum to a sending system of said
7	encrypted data block that identifies a data-stream that includes said data block
8	based on said read tag bits; and

- said receiving system decrypting said a payload data of said payload portion
   in response to receiving a decryption keys from said sending system.
  - 1 18. The method defined in claim 17 wherein said tag bits have a source identifier
  - 2 in said sending system of said decryption keys, and further including said receiving
  - 3 system sending said source identifier to said sending system in response to said
- 4 reading.
- 1 19. A data safeguarding system for a data block sent from a first system to a2 second system including:
  - a first system payload replacement circuit that replaces a portion of a payload of said data block with a tag data that indicates at least one decryption key for said data block in said first system;
  - a first system header flag setting circuit that sets a flag in a header of said data block when said data block includes said tag;
- 8 a first system encryption circuit that encrypts said payload with said keys; and
- 9 a first system data-stream sending circuit that sends a data-stream that
- 10 includes said data block to said second system after said header flag setting circuit
- 11 sets said flag and said encryption circuit encrypts said payload and said payload
- 12 replacement circuit replaces said portion of a payload.

- 1 20. The system defined in claim 19 wherein said first system encryption circuit
- 2 encrypts said portion of said payload.
- 1 21. The system defined in claim 19 further including at least one of
- a first system sending circuit that sends said at least one key to said second
- 3 system; and
- 4 a first system sending circuit that sends said at least one key and said portion
- 5 of said payload to said second system.
- 1 22. The system defined in claim 21 wherein said first system sending circuit
- 2 sending is based upon said first system receiving from said second system a first
- 3 datum that indicates at least one decryption key for said data block in said first
- 4 system
- 1 23. The system defined in claim 19, further including:
- 2 a second system header flag reading circuit that reads said flag in said
- 3 header;
- 4 a second system tag data reading circuit that reads said tag data if said
- 5 second system header flag reading circuit indicates that said flag includes said tag
- 6 data;
- 7 a second system data sending circuit that sends to said first system a datum
- 8 that identifies said data-stream based on said tag data; and
- a second system decrypting circuit that decrypts said encrypted block.

- 1 24. The system defined in claim 23 further including a first system key sending
- 2 circuit that sends said at least one key to said second system, and wherein said
- 3 second system decrypting circuit decrypts said data stream based on said at least
- 4 one key.
- 1 25. The system defined in claim 23 further including
- a first system sending circuit that sends said portion of said payload to said
- 3 second system in response to receiving from said second system a datum that
- 4 indicates said decryption keys in said first system
- 5 said first circuit encryption circuit further encrypts said replaced portion of
- 6 said payload;
- 7 a second system payload replacement circuit that replaces said received tag
- 8 data with said portion of said payload; and
- 9 said second system decrypting circuit further decrypts said portion of said
- 10 payload.
  - 1 26. The system defined in claim 19 wherein at least one of:
- 2 said sending occurs via a shared memory; and
- 3 said first system and said second system are separate physical devices; said
- 4 sending of said data-stream occurs on a first channel; and sending non-data-stream
- 5 data including at least one of said at least one key, said portion of said payload, and
- 6 said data-stream identifier occurs on a second channel.

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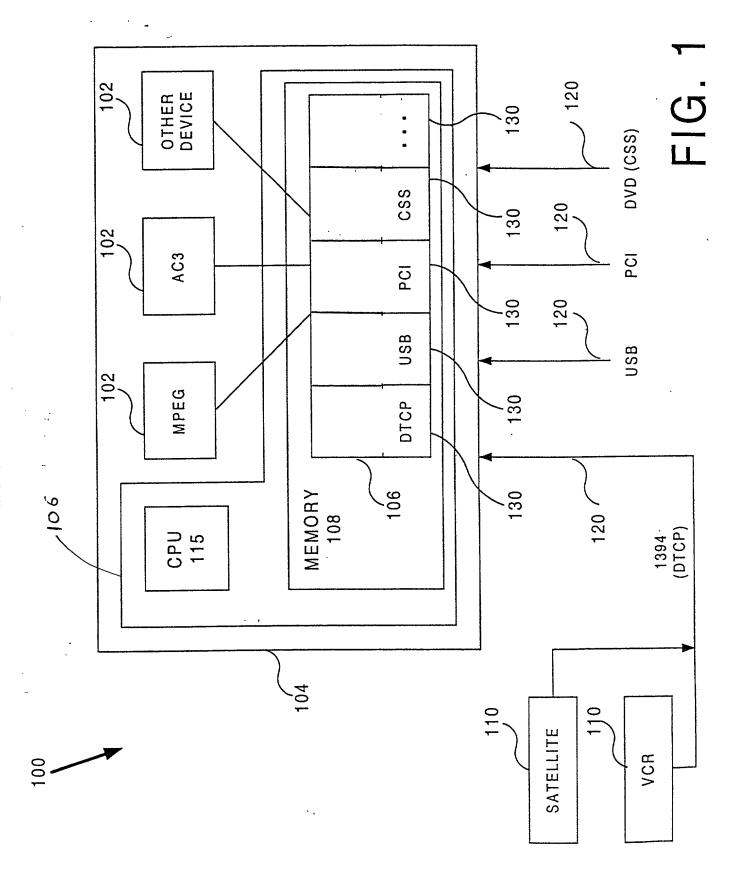
1	27. The system defined in claim 23 wherein said tag data further has an
2	identifier for accessing a first system unit that can send to said second system said
3	keys.

- 28. The system defined in claim 19 further including before said first circuit header flag setting circuit setting said flag and said first circuit encryption circuit encrypting said payload, a second circuit receiving circuit that can receive a stream of data from a third system wherein said data-stream is based on said stream of data.
  - 29. A system for safeguarding a data block of a data-streamsent from a first system to a second system comprising:
  - a second system header flag reading circuit that reads a flag in a header of said data block;
  - a second system tag data reading circuit that reads a data-stream identifier in a tag data of a payload portion of said block if said header flag reading circuit indicates that said flag includes said tag data; and
- a second system data sending circuit that sends to said first system a first datum that identifies said data-stream based on said data-stream identifier.
- 1 30. The system defined in claim 29 further including a second system 2 decrypting circuit that decrypts said data block.

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## ABSTRACT OF THE DISCLOSURE

A data safeguarding system, method, and article for safeguarding an encrypted data-stream transmitting on a first channel from a first system to a second system. The data-stream can be intertwined with other data-streams. The data-stream is arranged in fixed length sequential blocks, each block including a header portion and a payload portion. The first system places a flag marking in the header portion indicating that the payload includes a tag having at least one identifier for selecting the decryption keys from the first system. The second system reads the flag, and if the flag indicates a tag portion, reads the tag portion. The second system transmits the identifier to the first system on a second channel. The first system reads the identifier, retrieves the keys, and transmits the decryption keys to the second system on the second channel. The second system receives the decryption keys and decrypts the data block using the decryption keys.



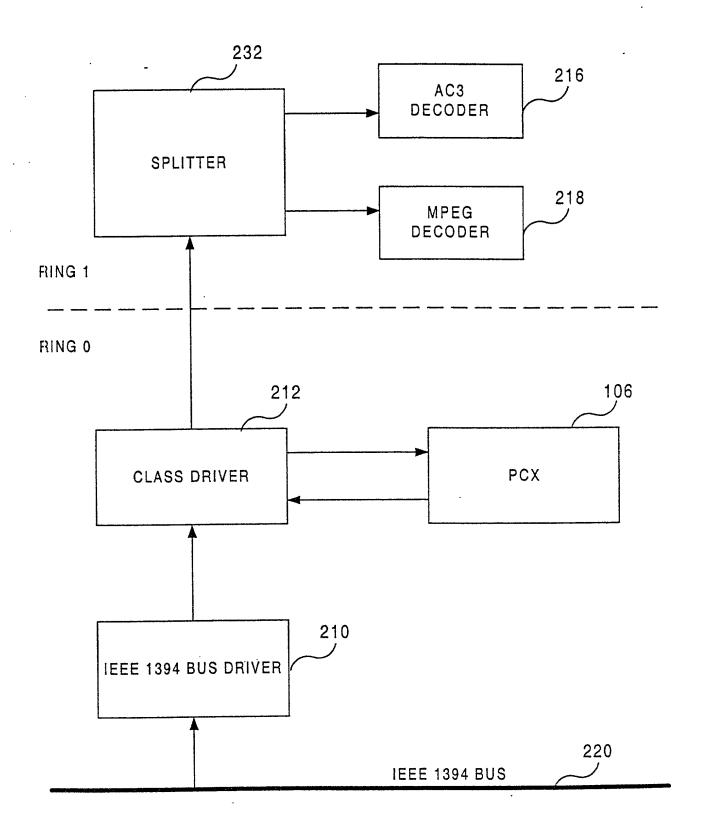
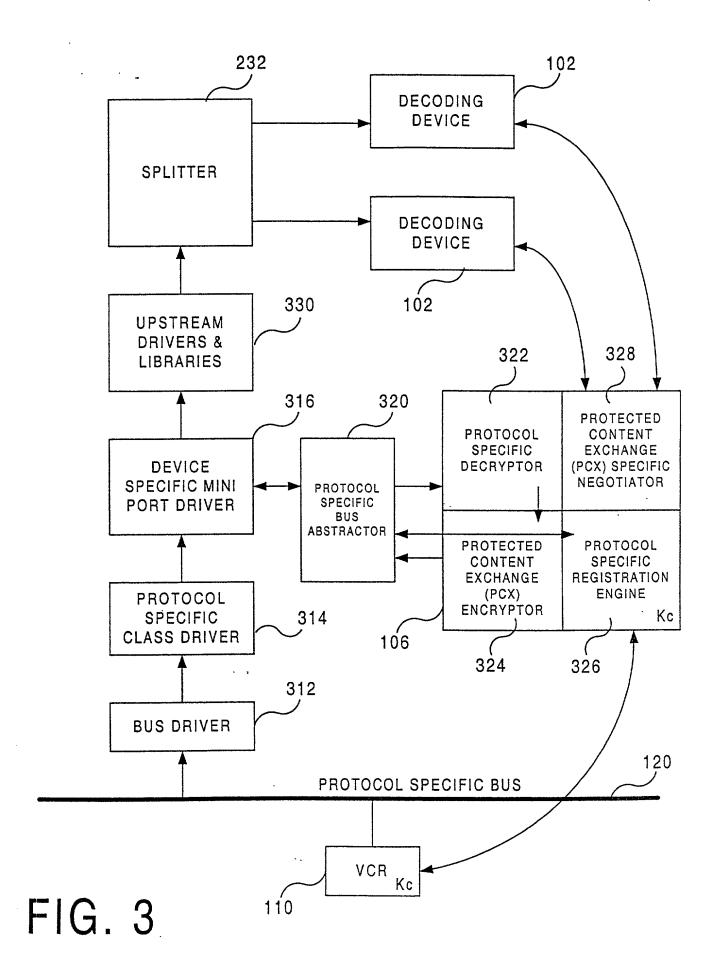
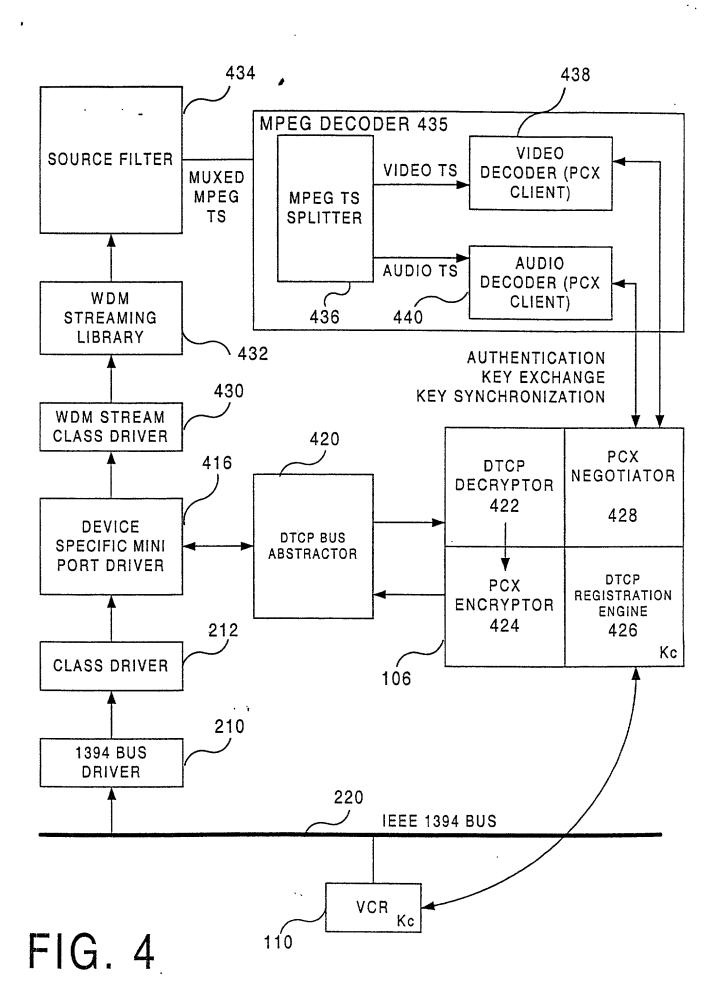
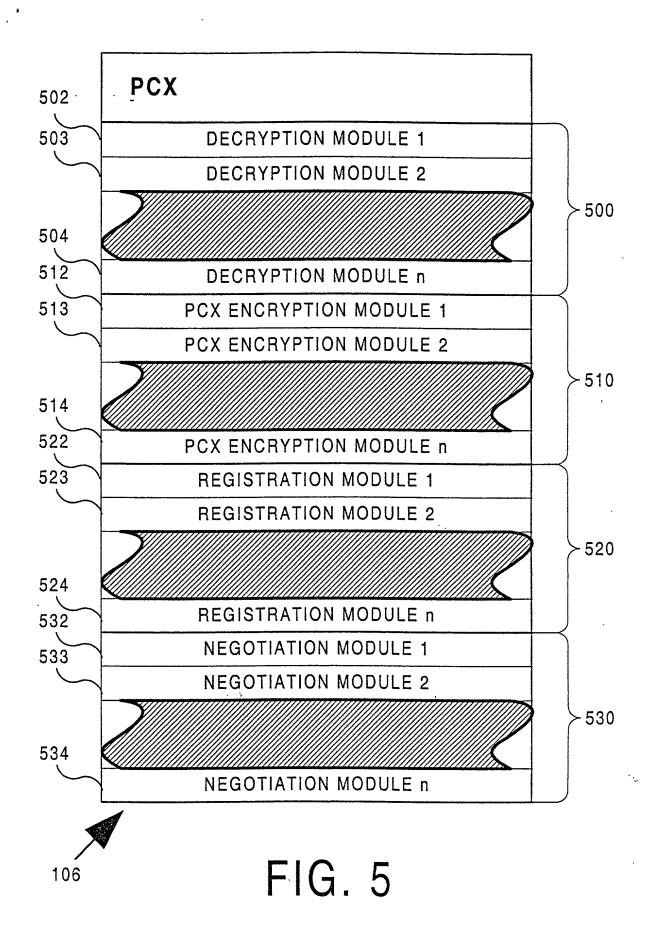
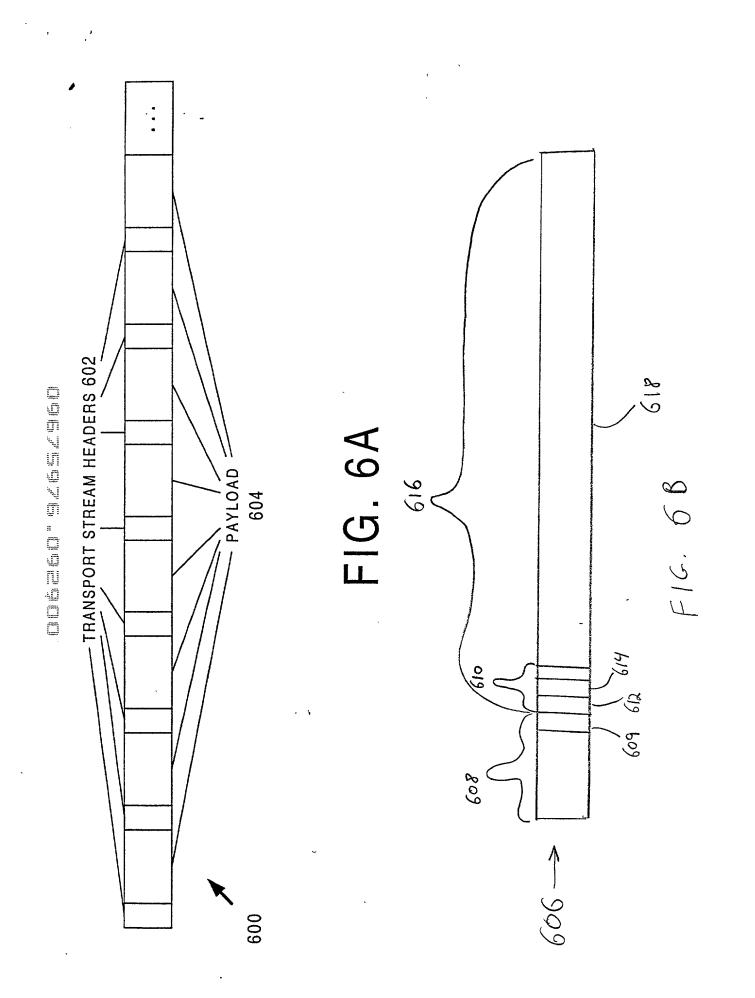


FIG. 2









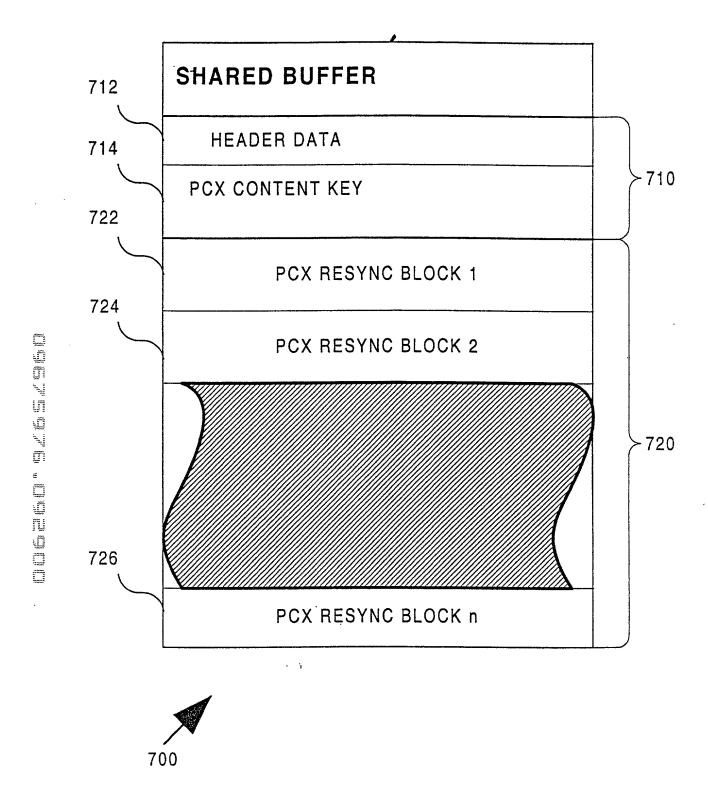
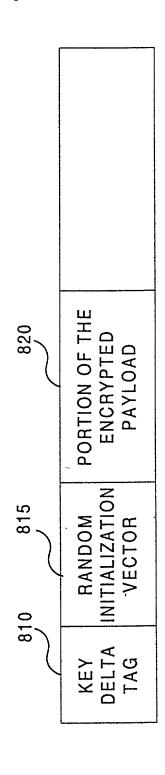


FIG. 7





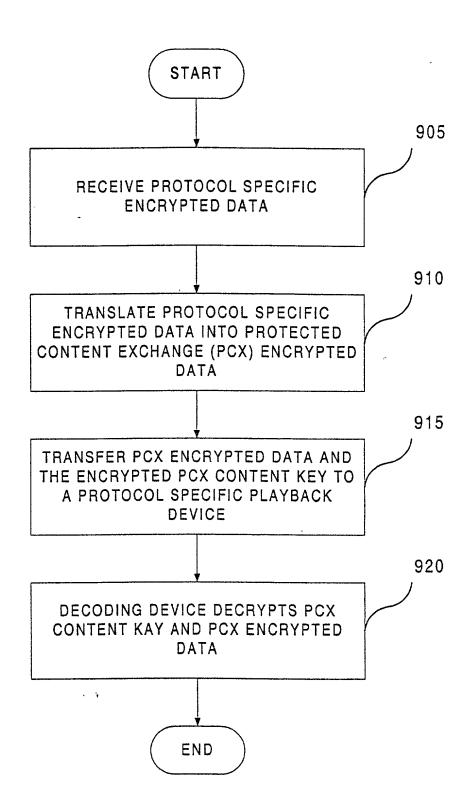


FIG. 9

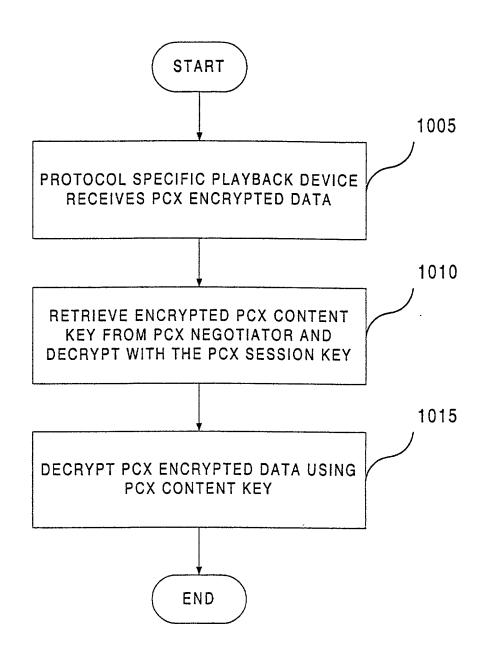
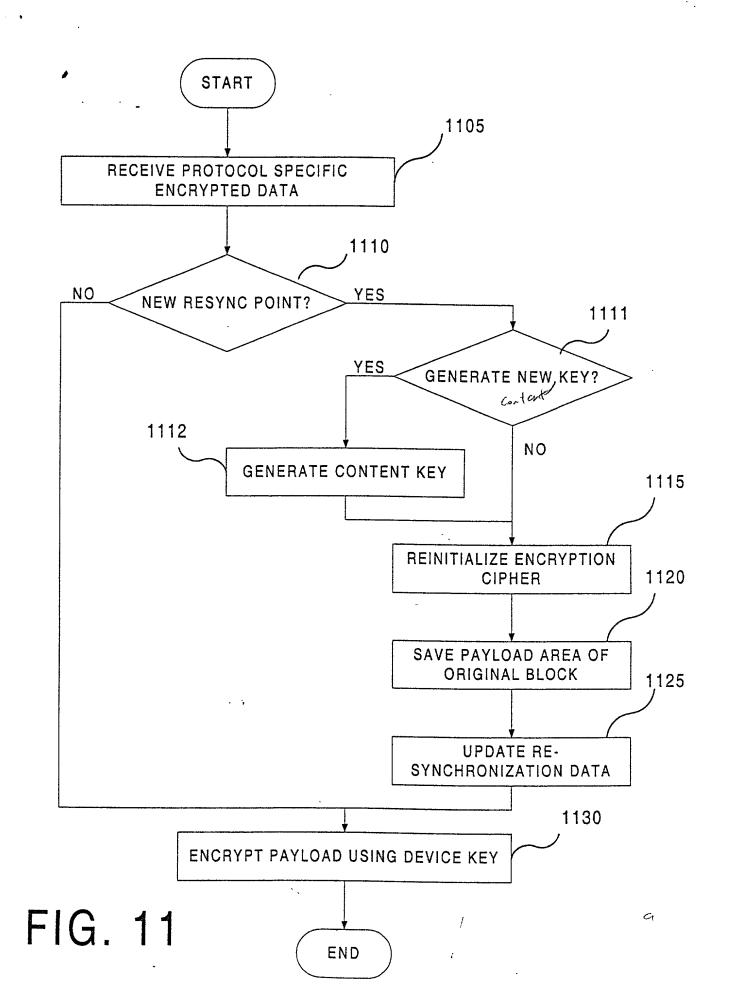
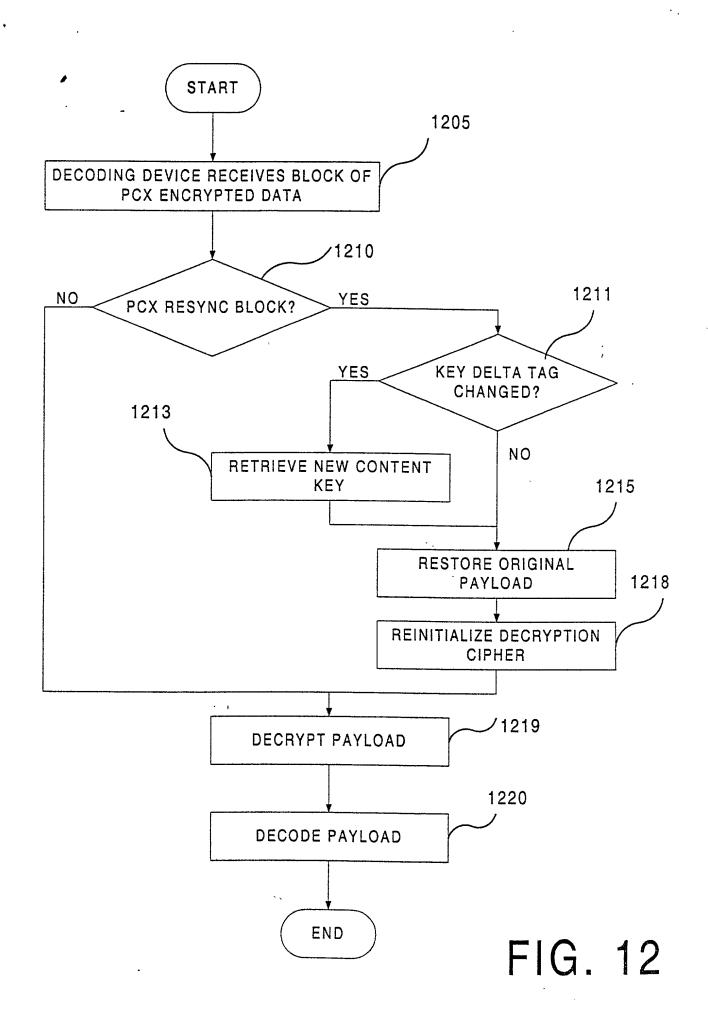


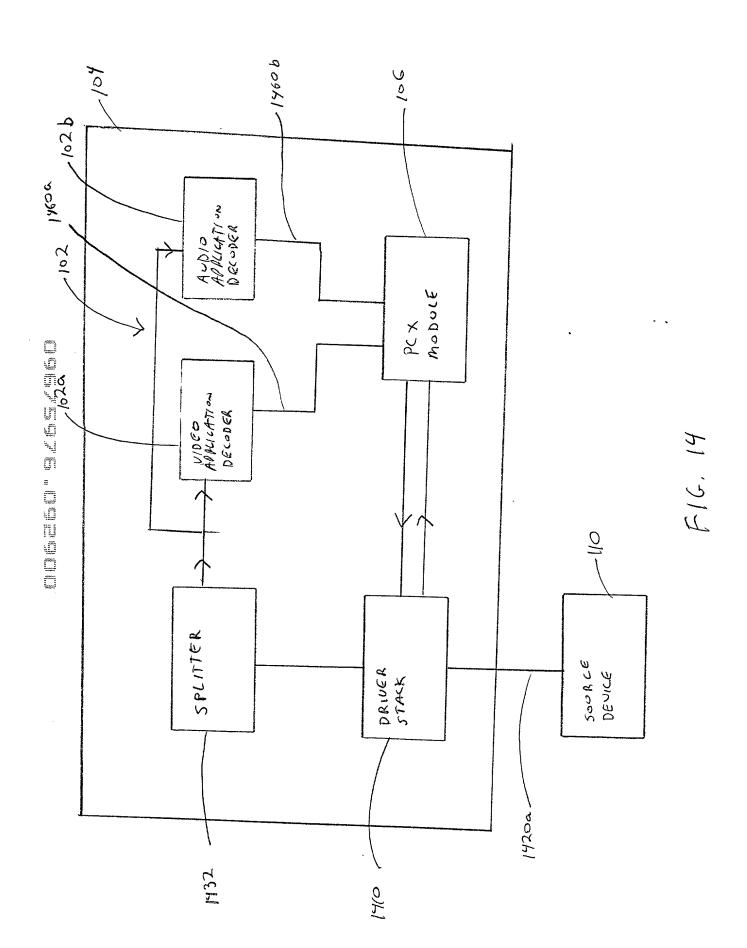
FIG. 10



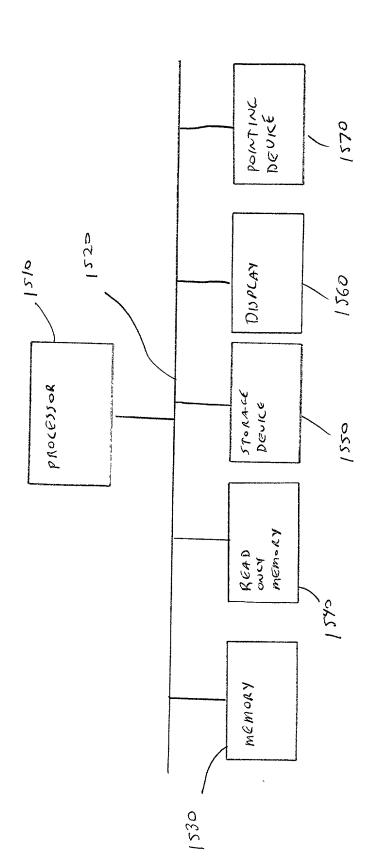


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FIG. 13

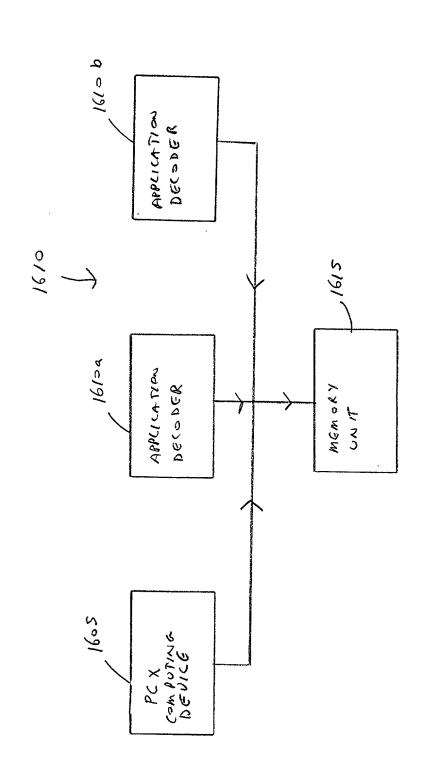


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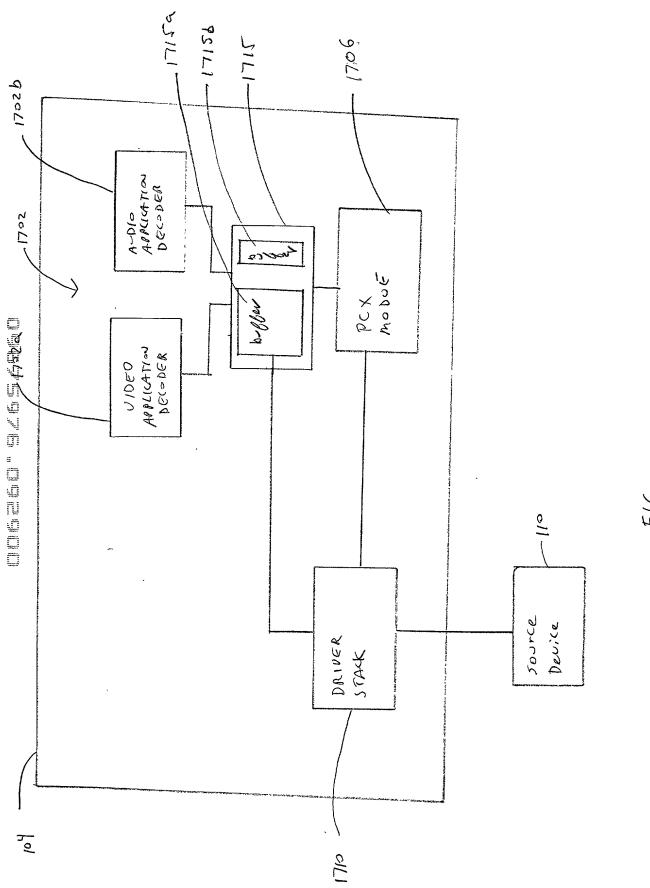


FIC. 15

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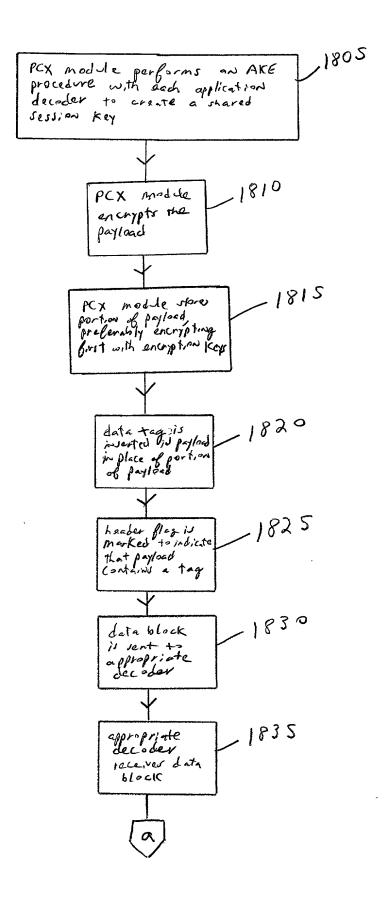


F16. 16



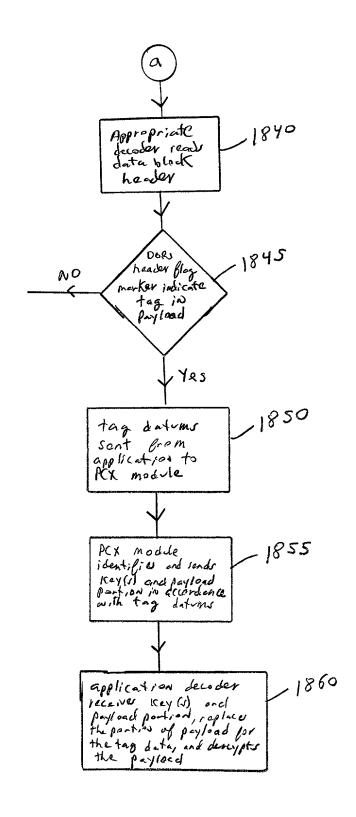
F16. 17

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F16.18

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F16:18 Continued

Attorney's Docket No.: 042390.P7957 **PATENT** 

# DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

As a below named inventor, I hereby declare that:					
My residence, post office address and citizenship are as stated below, next to my name.					
I believe I am the original, first, and sole inventor (if only one name is listed below) or an original, first, and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled  SYSTEM AND METHOD FOR SAFEGUARDING DATA BETWEEN A DEVICE DRIVER AND A DEVICE					
the specification of which					
or F		as Numberas cation Number(if applicable)			
I hereby state that I have respecification, including the content of the content	viewed and understand claim(s), as amended by	the contents of the above-ident y any amendment referred to ab	ified ove.		
I acknowledge the duty to di defined in Title 37, Code of	isclose all information k Federal Regulations, Se	nown to me to be material to par ection 1.56.	tentability	as /	
foreign application(s) for part	tent or inventor's certific atent or inventor's certif	5, United States Code, Section a ate listed below and have also in ficate having a filing date before	dentified	below	
Prior Foreign Application(s)			Priori <u>Claim</u>		
Number	Country	Day/Month/Year Filed	Yes	No	
Number	Country	Day/Month/Year Filed	Yes	No	
Number	Country	Day/Month/Year Filed	Yes	No	
I hereby claim the benefit under Title 35, United States Code, Section 119(e) of any United States provisional application(s) listed below:					
Application Number	Filing Date	<del></del>			
Application Number	Filing Date				

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, Section 112, I acknowledge the duty to disclose all information known to me to be material to patentability as defined in Title 37, Code of Federal Regulations, Section 1.56 which became available between the filing date of the prior application and the national or PCT international filing date of this application:						
Application Number	Filing Date	Status patented, pending, abandoned				
Application Number	Filing Date	Status patented, pending, abandoned				
I hereby appoint the persons listed on Appendix A hereto (which is incorporated by reference and a part of this document) as my respective patent attorneys and patent agents, with full power of substitution and revocation, to prosecute this application and to transact all business in the Patent and Trademark Office connected herewith.						
Send correspondence to K	enneth B. Paley	, BLAKELY, SOKOLOFF, TAYLOR &				
Send correspondence to <u>Kenneth B. Paley</u> , BLAKELY, SOKOLOFF, TAYLOR & (Name of Attorney or Agent)  ZAFMAN LLP, 12400 Wilshire Boulevard 7th Floor, Los Angeles, California 90025 and direct telephone calls to <u>Kenneth B. Paley</u> , (425) 827-8600.  (Name of Attorney or Agent)						
(Name of Attorney or Agent)  I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.						
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Post Office Address			
Full Name of Fifth/Jo	oint Inventor		
Inventor's Signature		Date	
Residence	(City, State)	Citizenship	(Country)
Post Office Address			
Full Name of Sixth/J	oint Inventor		
Inventor's Signature		Date	
Residence	(City, State)	Citizenship	(Country)
Post Office Address			
Full Name of Seven	th/Joint Inventor		
Inventor's Signature	)	Date	
Residence	(City, State)	Citizenship	(Country)
	3		

### APPENDIX A

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### APPENDIX B

#### Title 37, Code of Federal Regulations, Section 1.56 Duty to Disclose Information Material to Patentability

- (a) A patent by its very nature is affected with a public interest. The public interest is best served, and the most effective patent examination occurs when, at the time an application is being examined, the Office is aware of and evaluates the teachings of all information material to patentability. Each individual associated with the filing and prosecution of a patent application has a duty of candor and good faith in dealing with the Office, which includes a duty to disclose to the Office all information known to that individual to be material to patentability as defined in this section. The duty to disclosure information exists with respect to each pending claim until the claim is cancelled or withdrawn from consideration, or the application becomes abandoned. Information material to the patentability of a claim that is cancelled or withdrawn from consideration need not be submitted if the information is not material to the patentability of any claim remaining under consideration in the application. There is no duty to submit information which is not material to the patentability of any existing claim. The duty to disclosure all information known to be material to patentability of any claim issued in a patent was cited by the Office or submitted to the Office in the manner prescribed by §§1.97(b)-(d) and 1.98. However, no patent will be granted on an application in connection with which fraud on the Office was practiced or attempted or the duty of disclosure was violated through bad faith or intentional misconduct. The Office encourages applicants to carefully examine:
  - (1) Prior art cited in search reports of a foreign patent office in a counterpart application, and
- (2) The closest information over which individuals associated with the filing or prosecution of a patent application believe any pending claim patentably defines, to make sure that any material information contained therein is disclosed to the Office.
- (b) Under this section, information is material to patentability when it is not cumulative to information already of record or being made or record in the application, and
- (1) It establishes, by itself or in combination with other information, a prima facie case of unpatentability of a claim; or
  - (2) It refutes, or is inconsistent with, a position the applicant takes in:
  - (i) Opposing an argument of unpatentability relied on by the Office, or
  - (ii) Asserting an argument of patentability.

A prima facie case of unpatentability is established when the information compels a conclusion that a claim is unpatentable under the preponderance of evidence, burden-of-proof standard, giving each term in the claim its broadest reasonable construction consistent with the specification, and before any consideration is given to evidence which may be submitted in an attempt to establish a contrary conclusion of patentability.

- (c) Individuals associated with the filing or prosecution of a patent application within the meaning of this section are:
  - (1) Each inventor named in the application;
  - (2) Each attorney or agent who prepares or prosecutes the application; and
- (3) Every other person who is substantively involved in the preparation or prosecution of the application and who is associated with the inventor, with the assignee or with anyone to whom there is an obligation to assign the application.
- (d) Individuals other than the attorney, agent or inventor may comply with this section by disclosing information to the attorney, agent, or inventor.